

Algebraische Rechenmodelle, exercise sheet 10

January 26, 2015

Exercise 1

Let M be a BSS machine over \mathbb{R} and let $\alpha := (\alpha_1, \dots, \alpha_s) \in \mathbb{R}^s$ be the machine constants of M . Let $\phi_n^+(M)(\alpha, x_1, \dots, x_n)$ be a first order formula which is true iff M accepts $x \in \mathbb{R}^n$ and let $\phi_n^-(M)(\alpha, x)$ be a first order formula which is true iff M rejects $x \in \mathbb{R}^n$. Let $\tilde{\mathbb{R}}$ be a real closed extension of \mathbb{R} and define

$$\begin{aligned}x \in \tilde{A}_n &\Leftrightarrow \phi_n^+(M)(\alpha, x_1, \dots, x_n) \\x \in \text{comp}(\tilde{A}_n) &\Leftrightarrow \phi_n^-(M)(\alpha, x_1, \dots, x_n).\end{aligned}$$

Use the theorem below to show that \tilde{A}_n and $\text{comp}(\tilde{A}_n)$ are disjoint and that $\tilde{A}_n \cup \text{comp}(\tilde{A}_n) = \tilde{R}$. Show then that their definition does not depend on the machine M used to decide A over \mathbb{R} , i.e. any other machine deciding A over the reals will result in the same subsets of $(\tilde{R})^n$.

Theorem 1 (transfer principle). *Let (R, \leq) and $(\tilde{R}, \tilde{\leq})$ be real closed fields such that $R \subset \tilde{R}$ and $\tilde{\leq}$ extends the ordering \leq (i.e. $\forall x, y \in R$ we have $x \leq y$ iff $x \tilde{\leq} y$). Let ψ be a first order sentence (i.e. no free variables) which has only constants from R . Then ψ holds in $(\tilde{R}, \tilde{\leq})$ if and only if ψ holds in (R, \leq)*

Exercise 2

Let F be a subfield of the reals: $\mathbb{Q} \subseteq F \subseteq \mathbb{R}$. Call a set $A \subseteq \mathbb{R}$ *countably semi-algebraic over F* iff it is a countable union $A = \bigcup_{m \in \mathbb{N}} A_m$ of semi-algebraic sets A_m that are definable over F . Show that a language $L \subseteq \mathbb{R}^\infty$ is semi-decidable in the BSS model over \mathbb{R} if and only if L is countably semi-algebraic over some *finitely generated* field extension of \mathbb{Q} .

Exercise 3

Let $\tilde{\mathbb{R}}$ be a real closed extension of \mathbb{R} . Prove that $P_{\tilde{\mathbb{R}}} \neq NP_{\tilde{\mathbb{R}}} \Rightarrow P_{\mathbb{R}} \neq NP_{\mathbb{R}}$.

Hint: Use the fact that the problem (F^4, F_{zero}^4) over \mathbb{R} is $NP_{\mathbb{R}}$ -complete and the problem (F^4, F_{zero}^4) over $\tilde{\mathbb{R}}$ is $NP_{\tilde{\mathbb{R}}}$ -complete. Show with help of the transfer principle that a machine M deciding the problem (F^4, F_{zero}^4) over \mathbb{R} can be reinterpreted as a machine over \tilde{R} and then it decides the problem (F^4, F_{zero}^4) over $\tilde{\mathbb{R}}$ within the same time bound.